

# On the fixed points of the iterated pseudopalindromic closure

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# Iterated palindromic closure

## Palindromic closure

The *palindromic closure* of  $w \in \mathcal{A}^*$ , denoted  $w^+$ , is the shortest palindrome having  $w$  as prefix.

## Iterated palindromic closure

Let  $w \in \mathcal{A}^*$ . The iterated palindromic closure of  $w$ , denoted  $\text{Pal}(w)$ , is defined by

$$\text{Pal}(\varepsilon) = \varepsilon, \quad \text{Pal}(w) = (\text{Pal}(w[1, n-1]) \cdot w[n])^+.$$

## Example

$$\text{Pal}(123) = (\text{Pal}(12)3)^+ = ((\text{Pal}(1)2)^+3)^+ = ((12)^+3)^+ = ((121)3)^+ = 1213121.$$

# Generalization of the iterated palindromic closure

It can naturally be generalized to an infinite word  $u \in \mathcal{A}^\omega$  as

$$\text{IPal}(u) = \lim_{n \rightarrow \infty} \text{Pal}(u[1, n]).$$

## Standard Sturmian and standard episturmian words

Let  $w \in \mathcal{A}^\omega$ . Then  $\text{IPal}(w)$  is a *standard episturmian word* if  $|\mathcal{A}| \geq 3$  and if  $|\mathcal{A}| = 2$  and  $w \neq u\alpha^\omega$ , it is a *standard Sturmian word*.

## Example

Let consider the Tribonacci word  $T = \text{IPal}((123)^\omega)$ . Then,

$$T = \underline{1}2\underline{1}3\underline{1}21\underline{1}213121\underline{1}2131211213121 \dots$$

We denote by  $\Delta(T)$  the word that determines  $T$ , called the *directive word* of  $T$ . Here,  $\Delta(T) = (123)^\omega$ .

# Pseudopalindrome

An *antimorphism* of  $\mathcal{A}^*$  is a function  $\theta : \mathcal{A}^* \rightarrow \mathcal{A}^*$  such that for all  $u, v \in \mathcal{A}^*$ ,  
 $\theta(uv) = \theta(v)\theta(u)$ .

If  $\theta^2 = \text{id}$ , then it is *involutive*.

## $\theta$ -palindrome

For a fixed involutive antimorphism  $\theta$ , a finite word  $w \in \mathcal{A}^*$  is called a  *$\theta$ -palindrome (pseudopalindrome)* if  $\theta(w) = w$ .

In the sequel,  $R : \mathcal{A}^* \rightarrow \mathcal{A}^*$  is the involutive antimorphism defined as the reversal.

## Example

For  $w = abcab$ ,  $R(w) = bacba$ .

The  $R$ -palindromes are exactly the usual palindromes.

# Involutive antimorphisms

## Lemma

Let  $\tau$  be an involutive permutation over the alphabet  $\mathcal{A}$ . Then  $\theta = \tau \circ R = R \circ \tau$  is the unique involutive antimorphism on  $\mathcal{A}^*$  that extends  $\tau$ . Thus,

$$\theta(w) = \tau(w[n])\tau(w[n-1]) \cdots \tau(w[1]).$$

Any involutive antimorphism can be obtained that way.

## Example

Let  $\theta = R \circ \tau$ , with  $\tau(a) = b$  and  $\tau(b) = a$ . Then  $abab$  is a  $\theta$ -palindrome. Indeed,  $abab$  is a fixed point under  $\theta$  :

$$\theta(abab) = R(baba) = abab.$$

# Iterated pseudopalindromic closure

The  $\theta$ -palindromic (pseudopalindromic) closure of a word  $w \in \mathcal{A}^+$ , denoted  $w^\oplus$ , is the shortest  $\theta$ -palindrome having  $w$  as prefix.

## Example

Let  $\theta$  be an involutive antimorphism such that  $\tau(1) = 3$ ,  $\tau(2) = 4$ ,  $\tau(3) = 1$  and  $\tau(4) = 2$ , and let  $w = 123124$ . Then

$$w^\oplus = 1231 \cdot 24 \cdot \theta(1231) = 1231 \cdot 24 \cdot 3143.$$

The iterated pseudopalindromic closure, denoted  $\text{Pal}_\theta$ , is naturally defined by  $\text{Pal}_\theta(\varepsilon) = \varepsilon$ , and for  $w \in \mathcal{A}^*$ ,

$$\text{Pal}_\theta(w) = (\text{Pal}_\theta(w[1, n-1])w[n])^\oplus.$$

And  $\text{IPal}_\theta = \lim_{n \rightarrow \infty} \text{Pal}_\theta(w[1, n])$ .

# Questions

- 1 What do the fixed points of the iterated pseudopalindromic closure look like ?
- 2 How many are they ?
- 3 Do they have remarkable combinatorial properties ?



# Existence (1/2)

## Example

$$\text{IPal}_R(abx\cdots) = \underline{abax}\cdots$$

$$\text{IPal}_R^2(abx\cdots) = \underline{abaaba}\underline{x}\cdots$$

$$\text{IPal}_R^3(abx\cdots) = \underline{abaabaabaabaabaaba}\underline{abaaba}\underline{x}\cdots$$

Let  $E$  be the involutive antimorphism defined by  $E = R \circ \tau$ , with  $\tau(a) = b$  and  $\tau(b) = a$ . Then

$$\text{IPal}_E(abx\cdots) = \underline{abbaab}\underline{x}\cdots$$

$$\text{IPal}_E^2(abx\cdots) = \underline{abbaabbaabaabbaabbaabbaabbaabbaabbaabbaabbaab}\underline{x}\cdots$$

# Existence (2/2)

## Theorem and definition

Over a  $k$ -letter alphabet, with  $k \geq 2$ , there are 3 kinds of fixed points having at least 2 different letters, only depending on the first letters of the word and the involutory antimorphism  $\theta = R \circ \tau$  considered.

- 1 When  $\tau(a) = a$  and  $\tau(b) = b$ , with  $a \neq b$ , for a fixed  $n \geq 1$ ,  $\text{IPal}_\theta$  has a unique fixed point beginning with  $a^n b$ , denoted  $\mathbf{s}_{R,n,a,b}$ , which equals

$$\mathbf{s}_{R,n,a,b} = \lim_{i \rightarrow \infty} \text{Pal}^i(a^n b) = \underline{a}^n \underline{b} a^n (\underline{a} b a^n)^{n+1} \underline{b} (a^{n+1} b)^{n+1} a^n \underline{a} \cdots$$

- 2 When  $\tau(a) = a$  and  $\tau(b) = c$  for pairwise different letters  $a, b, c$ , for a fixed  $n \geq 1$ ,  $\text{IPal}_\theta$  has a unique fixed point beginning with  $a^n b$ , denoted by  $\mathbf{s}_{\mathcal{H},n,a,b,c}$ , which equals

$$\mathbf{s}_{\mathcal{H},n,a,b,c} = \lim_{i \rightarrow \infty} \text{Pal}_{\mathcal{H}}^i(a^n b) = \underline{a}^n \underline{b} c a^n \underline{c} b a^n b c a^n (\underline{a} b c a^n c b a^n b c a^n)^n \underline{c} \cdots$$

- 3 When  $\tau(a) = b$  and  $\tau(b) = a$ , with  $a \neq b$ ,  $\text{IPal}_\theta$  has a fixed point beginning with  $a^n b$  only if  $n = 1$ . It is denoted by  $\mathbf{s}_{E,a,b}$  and equals

$$\mathbf{s}_{E,a,b} = \lim_{i \rightarrow \infty} \text{Pal}_E^i(a) = \underline{a} b \underline{b} a \underline{a} b \underline{b} a \underline{a} b \underline{b} a \underline{a} b \underline{b} a \underline{a} b \underline{b} a \underline{a} b \underline{b} \cdots$$

# Combinatorial properties of $\mathbf{s}_{R,n,a,b}$

- 1 For a fixed positive  $n \in \mathbb{N}$ ,  $\mathbf{s}_{R,n,a,b}$  is not ultimately periodic and consequently, is standard Sturmian (using Droubay, Justin, Pirillo - 2001).
- 2 For a fixed  $n \in \mathbb{N}$ ,  $\mathbf{s}_{R,n,a,b}$  is not a fixed point of a nontrivial morphism (using Arnoux, Rauzy - 1991 and Crisp et al - 1993).
- 3  $\mathbf{s}_{R,n,a,b}$  is  $(n+4)$ -th power-free, but contains  $(n+3)$ -th powers (using Vandeth - 2000).
- 4 For any  $n \geq 1$ ,  $\alpha_{n,a,b}$  is transcendental (using Adamczewski, Bugeaud - 2007).

# Combinatorial properties of $\mathbf{s}_{E,a,b}$ (1/2)

## Useful property [de Luca, De Luca - 2006]

Let  $\theta = \tau \circ R$  be an involutory antimorphism over an alphabet  $\mathcal{A}$ , with  $\mu_\theta$  the morphism defined for all  $a$  in  $\mathcal{A}$ , by  $\mu(a) = a$  if  $a = \tau(a)$  and by  $\mu(a) = a\tau(a)$  otherwise.

Then, for any  $\mathbf{w} \in \mathcal{A}^\omega$  and for any involutory antimorphism  $\theta$ , one has

$$\text{IPal}_\theta(\mathbf{w}) = \mu_\theta(\text{IPal}(\mathbf{w})).$$

## Idea

First consider  $\mathbf{w}_E = \text{IPal}(\mathbf{s}_{E,a,b})$  and then, extend its properties to  $\mu_E(\text{IPal}(\mathbf{s}_{E,a,b})) = \mathbf{s}_{E,a,b}$ .

$$\mathbf{w}_E = \underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b\underline{a}b \dots$$

# Combinatorial properties of $\mathbf{s}_{E,a,b}$ (2/2)

## Property of $\mathbf{w}_E$

$\mathbf{w}_E$  is not ultimately periodic, and consequently, is a Sturmian word.

## Lemma

An infinite word  $\mathbf{w}$  over  $\mathcal{A}$  is ultimately periodic if and only if  $\mu_\theta(\mathbf{w})$  is so.

$\implies \mathbf{s}_{E,a,b}$  is not ultimately periodic.

## Lemma

$\mathbf{w}_E$  is not a fixed point for some non-trivial morphism  $\implies \mathbf{s}_{E,a,b}$  is not a fixed point for some non-trivial morphism.

## Proposition

$\mathbf{w}_E$  and  $\mathbf{s}_{E,a,b}$  both contain 4-th powers, but no 5-th power words (using Shur - 2000).

# Combinatorial properties of $\mathbf{s}_{\mathcal{H},n,a,b,c}$ (1/3)

Recall :  $\tau(a) = a$ ,  $\tau(b) = c$  and  $\tau(c) = b$ , and  $\theta = R \circ \tau$ .

First consider  $\mathbf{w}_{\mathcal{H}} = \text{IPal}(\mathbf{s}_{\mathcal{H},n,a,b,c})$  and then extend its properties to  $\mu_{\mathcal{H}}(\text{IPal}(\mathbf{s}_{\mathcal{H},n,a,b,c})) = \mathbf{s}_{\mathcal{H},n,a,b,c}$ .

$$\mathbf{w}_{\mathcal{H},n} = \text{IPal}(\mathbf{s}_{\mathcal{H},n,a,b,c}) = \underline{a}^n \underline{b} a^n \underline{c} a^n b a^n \underline{a} b a^n c a^n b a^n \dots$$

## Property of $\mathbf{w}_{\mathcal{H},n}$

$\mathbf{w}_{\mathcal{H},n}$  is not ultimately periodic, and consequently, is a strict standard episturmian word.

$\implies \mathbf{s}_{\mathcal{H},n,a,b,c}$  is not ultimately periodic.

## Proposition

$\mathbf{w}_{\mathcal{H},n}$  is not a fixed point of a nontrivial morphism (using Justin, Pirillo - 2002).

## Combinatorial properties of $\mathbf{s}_{\mathcal{H},n,a,b,c}$ (2/3)

Generalization of a result of Justin, Pirillo - 2002 for strict episturmian word having periodic directive word :

### Proposition

Let  $\mathbf{s}$  be a strict standard episturmian word directed by a word  $\Delta$  and let  $\ell$  denotes the greatest integer such that  $\alpha^\ell$  is a factor of  $\Delta$  with  $\alpha$  a letter. Assume  $\Delta$  contains at least one factor  $aua^\ell va$  with  $a$  a letter and  $u, v$  non empty words that do not contain the letter  $a$ . Then  $\mathbf{s}$  is  $(\ell + 3)$ -th power-free but contains an  $(\ell + 2)$ -th power.

### Corollary

The words  $\mathbf{w}_{\mathcal{H},n}$  are  $(n + 4)$ -th power free but contain  $(n + 3)$ -th powers.

# Combinatorial properties of $\mathbf{s}_{\mathcal{H},n,a,b,c}$ (3/3)

Let  $\mathbf{s}_{\mathcal{H},n,a,b,c}$  be a fixed point of the  $\text{IPal}_{\mathcal{H}}$  operator, for a fixed  $n$ . Then  $\mathbf{s}_{\mathcal{H},n,a,b,c}$  satisfies the following properties :

- 1 It is not an episturmian word, but is a pseudostandard word.
- 2 It is not a fixed point for some non trivial morphism.
- 3 It is  $(n + 4)$ -th power-free but contains  $(n + 3)$ -th powers.
- 4 The frequencies of the letters  $b$  and  $c$  are equal.



# Open problems

- What is the critical exponent : its value, the number of occurrences, their positions, etc.
- Does there exist a connection between  $s_{R,n,a,b}$  and  $s_{R,n+1,a,b}$  ? between  $s_{\mathcal{H},n,a,b,c}$  and  $s_{\mathcal{H},n+1,a,b,c}$  ?
- Can we give a geometric interpretation of the iterated palindromic (pseudopalindromic) closure ?
- What are the letter frequencies of  $s_{\mathcal{H},n,a,b,c}$  ?
- Any remarkable combinatorial properties ?